

# Architectural Patterns Revisited – A Pattern Language

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Architectural patterns are a key concept in the field of software architecture: they offer well-established solutions to architectural problems, help to document the architectural design decisions, facilitate communication between stakeholders through a common vocabulary, and describe the quality attributes of a software system as forces. Regrettably, finding and applying the appropriate architectural patterns in practice still remains largely ad-hoc and unsystematic. This is due to the lack of consensus in the community with respect to the “philosophy” and granularity of architectural patterns, as well as the lack of a coherent pattern language. In this paper we attempt to establish common ground in the architectural patterns community by proposing a pattern language that acts as a superset of the existing architectural pattern collections and categorizations. This language is particularly focused on establishing the relationships between the patterns and performs a categorization based on the concept of “architectural views”.

## 1 Motivation

The software architecture community has had many debates on various aspects like which design and evaluation methods, Architecture Description Languages, and views are best for which case. Architectural patterns were one of the few points, where consensus was achieved in the field of software architecture: their significance is well-established and they are essential to an architecture description [SC97, SG96, BCK98, CBB<sup>+</sup>02, BMR<sup>+</sup>96, SSRB00]. Regrettably, describing, finding, and applying architectural patterns in practice still remains largely ad-hoc and unsystematic. This is due to several issues that are still not resolved: there is a semantic gap concerning what architectural patterns really represent, what is the philosophy behind them; moreover, there is much confusion with respect to what is the granularity of architectural patterns; finally there is no accepted classification or cataloguing of patterns that can be used by architects. The next paragraphs elaborate on these issues.

There are two different “schools of thought” in the literature with respect to the nature of architectural patterns: one that uses the term “architectural pattern” (e.g. in [BMR<sup>+</sup>96, SSRB00, VKZ04]) and another that uses the term “architectural style” (e.g. in [SG96, SC97, BCK98, CBB<sup>+</sup>02]). Both terms refer to recurring solutions that solve problems at the architectural design level, and provide a common vocabulary in order to facilitate communication. They both

also accept that patterns provide the means to reason for the quality attributes of a software system and help to document the design decisions taken by the architect. But they have some key differences in their underlying philosophy:

- In the architectural patterns perspective, patterns are considered as problem-solution pairs that occur in a given context and are affected by it. Furthermore a pattern does not only document “how” a solution solves a problem but also “why” it is solved, i.e. the rationale behind this particular solution. In particular, the description of the problem pays much attention to the forces that shape the problem, while the solution elaborates on how (and if) those forces are resolved. The description of architectural patterns is based on the context-problem-solution triplet and may be further elaborated with richer details, especially focusing on the rationale behind a solution. Moreover, patterns are meant to work synergistically in the context of a pattern language, and have numerous inter-dependencies among each other. Finally, there are a number of postulations for a solution to qualify as a pattern. For instance, the pattern must capture common practice (e.g. have at least three known uses) and at the same time the solution of the pattern must be non-obvious. Patterns should provide aesthetic solutions, and in the pattern literature the human aspect of software is accentuated [Cop96].
- In the architectural styles perspective, the problem does not receive much attention nor does the rationale behind choosing a specific solution. In [SC97, BCK98] a style is looked upon in terms of components, connectors, and issues related to control and data flow. In [SG96] attention is drawn to architectural configurations, semantics of the styles, and potential architectural analysis that can be performed on the systems built on the styles. In [BCK98, CBB<sup>+</sup>02] the concept of architectural styles is treated as a set of constraints on components, connectors, and their interactions. Similarly, in [MM03], an architectural style is represented by components, connectors, their configurations, and constraints upon all of them. In this viewpoint, patterns are not considered generic and “timeless” in the Alexandrian sense [AIS<sup>+</sup>77, Ale79], but become much more concrete and focused. This also leads to multiple variations of the same pattern in order to solve specialized design problems [SC97]. Also, since the pattern’s implementation details can be pinpointed, Architecture Description Languages [MT00] can be designed in order to support individual architectural patterns.

As far as the granularity of architectural patterns is concerned, it is usually not clear when a pattern is “big” enough to be considered architectural. In particular, the category of “design patterns”, e.g. as described in [GHJV94], are often referred to, or used as architectural patterns. In general it is hard to draw the line between architectural patterns and design patterns. In fact, it depends heavily on the viewpoint of the designer or architect whether a specific pattern is categorized as an architectural pattern or a design pattern. Consider for instance, a classical design pattern, the INTERPRETER pattern [GHJV94]. The description in [GHJV94] presents it as a concrete design guideline. Yet, instances of the pattern are often seen as a central elements in the architecture of software systems, because an INTERPRETER is a central, externally visible component – i.e. the pattern is treated like an architectural pattern (see [SG96]).

Finally, there is no single catalogue of architectural patterns for software architects to use. Instead there is a voluminous and heterogeneous literature of patterns, where the various patterns

differ in their philosophy and way of description and are often not related in the context of a pattern language. To make matters worse, many architectural patterns languages have been developed since the earlier software patterns literature [SG96, SC97, BMR<sup>+</sup>96, GHJV94] has been documented, but the former are not clearly related to the latter. Of course, there have been attempts to classify architectural patterns: in [BMR<sup>+</sup>96] architectural patterns are categorized according to the global system properties that they support; in [SC97, BCK98] architectural patterns are classified with respect to a framework of features, like the types of components and connectors, and control and data issues; a more recent classification scheme that has been proposed is based on the concept of architectural views [CBB<sup>+</sup>02]. But again there is no consensus on these classifications that could possibly lead to a single scheme.

## 2 Putting the pieces together: a pattern language

### 2.1 The approach

In this paper we attempt to find *common ground* in the architectural patterns community by proposing a pattern language that acts as a superset of the existing architectural pattern collections and categorizations. This language, as a whole, is greater than the sum of its parts because it particularly focuses on establishing the relationships between the patterns in order to present the “big picture”. In particular this pattern language tackles the aforementioned shortcomings in the following ways:

- We consider that both architectural patterns and architectural styles are in essence the same concepts and that they only differ in using different description forms. Therefore, we put both the “classical” architectural patterns such as those from POSA (see [BMR<sup>+</sup>96]) and the architectural styles from SEI (see [SC97, Sha96, SG96, BCK98, CBB<sup>+</sup>02]) in the same pattern language, paying special attention on relating them with each other. For the sake of simplicity, we shall use only the term “architectural pattern” for the rest of this paper. Note that we are aware that some patterns described below, e.g. EXPLICIT INVOCATION would not qualify as patterns in a strict interpretation of the pattern definition, because they may be considered as stating really obvious solutions. Naturally, the obviousness of a solution is a matter of definition. We nevertheless include them in our pattern language because they are important pieces of the architectural patterns puzzle and substantially contribute in putting the individual architectural styles and patterns together into a coherent pattern language.
- We consider a pattern to be architectural, if it refers to a problem at the architectural level of abstraction; that is, if the pattern covers the overall system structure and not only a few individual subsystems. Thus we have not included the classical “design patterns” (e.g. [GHJV94]) in our pattern language except for one (INTERPRETER). However we emphasize that these design patterns can potentially be used as architectural patterns, if one applies them at the level and scope of a system’s architecture.
- We propose a classification of the architectural patterns, based upon architectural views, that extends the one proposed in [CBB<sup>+</sup>02]. The next subsection elaborates on this classification scheme.

This pattern language, as aforementioned, contains patterns from existing collections of architectural patterns. The emphasis of this language is therefore not on describing the individual patterns; they have already been elaborately described before. For space restrictions we don't repeat these descriptions here. Instead emphasis is given only on the related pattern sections that analytically describe the relationships between the patterns. Consequently, the intended audience for this pattern language is comprised of architects that have a sound knowledge of these patterns and can accordingly understand the “big picture” of the inter-related patterns.

## 2.2 Classification of architectural patterns according to views

The classification scheme for architectural patterns that we propose is based on the concept of architectural views.

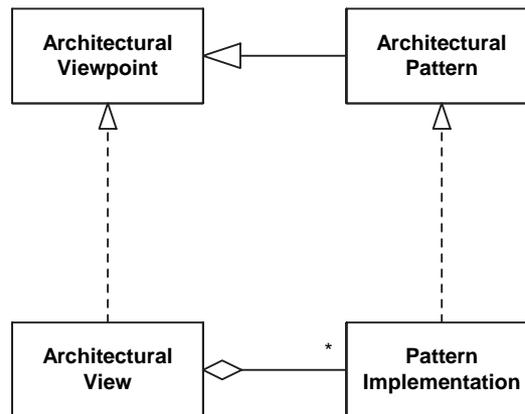


Figure 1: Architectural views, viewpoints, and patterns

An **Architectural View**<sup>1</sup> is a representation of a system from the perspective of a related set of *concerns* [IEE00] (e.g. a concern in a distributed system is how the software components are allocated to network nodes). This representation is comprised of a set of system *elements* and the *relationships* associated with them [CBB<sup>+</sup>02]. The *types* of the elements and the relationships as well as other meta-information on the views are described by the **Viewpoint** [IEE00] in order to document and communicate views unambiguously. Therefore a view is an instance of a viewpoint for a particular system, because the elements and relationships contained in the View, are instances of the corresponding types of elements and relationships contained in the Viewpoint.

An **Architectural Pattern**, on the other hand, also defines types of elements and relationships that work together in order to solve a particular problem from some perspective. In fact, an architectural pattern can be considered as a specialization of a viewpoint since it provides specialized semantics to the types of elements and relationships, as well as constraints upon them.

<sup>1</sup>We use the term *view* to describe the different perspectives of the architecture like “structural view” or “behavioral view”, instead of “structural architecture” or “behavioral architecture”, in the spirit of [IEE00] which mandates that a system has *one* architecture, whose description can be organized in many views.

An essential issue however is how architectural views relate to architectural patterns. There have been two major approaches on this matter so far:

- The first considers that views are of large granularity in the sense that the elements and relationships are generically defined. Therefore, in such a coarse-grained view, multiple architectural patterns may be applied. For instance we can consider a *structural* view, that describes how a system is structurally decomposed into components and connectors. In this view, we may apply the LAYERS and the PIPES AND FILTERS patterns. Some well-established examples of this approach are Kruchten's "4+1 views" [Kru95], the so-called "Siemens 4 views model" [HNS00], and the "Zachman framework" [Zac87]. The same thesis is supported in the IEEE 1471 Recommended Practice for Architectural Description. Even though it does not prescribe any specific set of views, it considers views of such granularity, e.g. structural and behavioral views.
- The second considers that each architectural pattern corresponds to a view in a one-to-one-mapping, and is mainly advocated in [CBB<sup>+</sup>02]. This notion of a view is of course far more fine-grained since it leads to elements and relationships of very specialized semantics, e.g. consider a pipe-and-filter view or a client-server view. These views are categorized into classes, called **Viewtypes**, which are of the same granularity as the views of the first approach. For example, in [CBB<sup>+</sup>02] the Components and Connectors viewtype contains the pipe-and-filter and the client-server views.

We follow the middle path between the two aforementioned approaches. We consider that views should be more fine-grained than the first approach in order to be useful. Therefore, instead of generic structural or behavioral views, we consider views that show more specific aspects of the system like the flow of data or the interaction of components. On the other hand we consider views to be more coarse-grained than individual architectural patterns, since more than one pattern can be either complimentary or alternatively applied in a given view. For example in a view that shows how shared data is manipulated by a number of components, we could either apply SHARED REPOSITORY or ACTIVE REPOSITORY.

Figure 1 illustrates the relationships between views, viewpoints, and patterns, using the Unified Modeling Language (UML). A viewpoint is related with a realization dependency with a view so that, according to the UML semantics, the viewpoint represents a specification and a view represents an implementation of the former. Our classification scheme for architectural patterns is organized around a number of the most common views that are used in practice. Therefore an architectural pattern is classified in a particular view if the pattern implementation is part of this view; or in other words if the pattern is a specialization of the viewpoint that types the particular view. Emphasis is given on the multiple and complex inter-relationships between the patterns.

In our classification scheme each pattern is assigned to one primary view, which is the most suitable. However there are some cases where a single pattern could be used in a second or third view. This can be derived as follows: when two patterns from different views are combined on the same system, then they can be seen in both views. For example, consider the case of a SHARED REPOSITORY which also implements a CLIENT-SERVER structure because the repository plays the role of a server satisfying the clients-accessors requests. In this case each of the two patterns is visible both in the data-centered and the component interaction views.

The views that we have chosen for this classification scheme contain mainly two types of elements: *components* which are units of runtime computation or data-storage, and *connectors* which are the interaction mechanisms between components [PW92, CBB<sup>+</sup>02]. There are of course other views that contain different kinds of elements, which we did not include in our pattern language for the time being. For instance in [CBB<sup>+</sup>02] the “module” views deal with implementation modules (e.g. Java or C++ classes), while the “allocation” views deal with how software elements are allocated to environment elements (e.g. code units are allocated to members of the development team).

### 2.3 Overview of the pattern language

The following pattern language is comprised of component and connector views, and the patterns that are classified in each view. Note that we focus on “classical” architectural patterns from both POSA (see [BMR<sup>+</sup>96]) and SEI (see [SC97, Sha96, SG96, BCK98, CBB<sup>+</sup>02]), because these have been well-established in the software architecture community. We have also included a few patterns from other sources, in order to describe important links or gaps in the realm of the “classical” architectural patterns. We also discuss some links to related pattern languages inside the description of the patterns.

- The Layered View deals with how the system as a complex heterogeneous entity can be decomposed into interacting parts.
  - LAYERS [SC97, Sha96, SG96, BCK98, CBB<sup>+</sup>02, BMR<sup>+</sup>96]
  - INDIRECTION LAYER [Zdu04, Zdu03] (a variant of this pattern is called “virtual machine” in [CBB<sup>+</sup>02])
- The Data Flow View deals with how streams of data are successively processed or transformed by components.
  - BATCH SEQUENTIAL [SG96, SC97, BCK98]
  - PIPES AND FILTERS [SC97, Sha96, SG96, BCK98, CBB<sup>+</sup>02, BMR<sup>+</sup>96]
- The Data-centered View is appropriate when the concerns involve how a central repository of data is accessed by multiple components.
  - SHARED REPOSITORY [VKZ04] (called “repository” in [SG96, Sha96, BCK98])
  - ACTIVE REPOSITORY (called “blackboard” in [SC97, BCK98])
  - BLACKBOARD [BMR<sup>+</sup>96, SG96]
- The Adaptation View deals with how the system adapts itself during evolution.
  - MICROKERNEL [BMR<sup>+</sup>96]
  - REFLECTION [BMR<sup>+</sup>96]
  - INTERCEPTOR [SSRB00]
- The Language Extension View is concerned with how systems offer an abstraction layer to the computation infrastructure.

- INTERPRETER [SG96, Sha96, BCK98, GHJV94]
- VIRTUAL MACHINE [GMSM00]
- RULE-BASED SYSTEM [SG96, BCK98]
- The User Interaction View shows the runtime structure of components that offer a user interface.
  - MODEL-VIEW-CONTROLLER [KP88, BMR<sup>+</sup>96]
  - PRESENTATION-ABSTRACTION-CONTROL [Cou87, BMR<sup>+</sup>96]
  - C2 [TMA<sup>+</sup>96]
- The Component Interaction View focuses on how individual components exchange messages but retain their autonomy.
  - EXPLICIT INVOCATION (called “communicating processes” in [SC97, Sha96, BCK98, CBB<sup>+</sup>02])
  - IMPLICIT INVOCATION [Sha96, BCK98, SG96] (also called “event systems” in [SC97, SG96, BCK98])
  - CLIENT-SERVER [SC97, SG96, BCK98, CBB<sup>+</sup>02]
  - PEER-TO-PEER [CBB<sup>+</sup>02]
  - PUBLISH-SUBSCRIBE [BCK98, CBB<sup>+</sup>02] (called “publisher-subscriber” in [BMR<sup>+</sup>96]).
- The Distribution View tackles concerns about disseminating components in a networked environment.
  - BROKER [BMR<sup>+</sup>96, VKZ04]
  - REMOTE PROCEDURE CALLS [VKZ04] (called “distributed objects” in [SC97])
  - MESSAGE QUEUEING (called “messaging” in [HW03, VKZ04])

The next sections elaborate on the various views and the patterns assigned to each view. The description of each view, i.e. the *viewpoint*, is presented informally with respect to the types of elements and relationships, as well as the concerns addressed by the view. Each pattern is described using a brief summary, and a more elaborate discussion of its relationships to other patterns.

### 3 Layered View

In the Layered View the system is viewed as a complex heterogeneous entity that can be decomposed into interacting parts. The concerns addressed by this view are:

- What are the parts that make up the whole system?
- How do these parts interact with each other?
- How do the parts perform their functionality and still remain decoupled from each other?

- How are the quality attributes of modifiability, scalability, and integrability supported?

The individual parts of the system are components that are decoupled as much as possible from one another. The interaction mechanisms between the components are implemented through connectors that include appropriate interfaces, states, and interaction protocols. There is usually an overall control mechanism that maintains an overall organization scheme by orchestrating the various components.

Figure 2 illustrates patterns and their relationships from the Layered, Data Flow, and Data-centered Views.

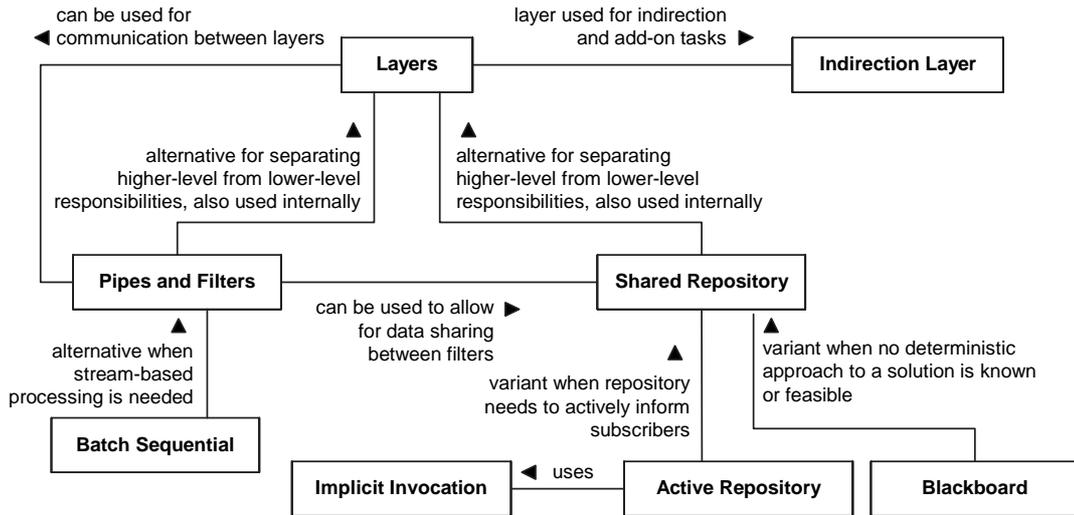


Figure 2: Overview: Patterns of the Layered, Data Flow and Data-centered Views

### Pattern: Layers

Consider a system in which high-level components depend on low-level components to perform their functionality, which further depend on even lower-level components and so on. Decoupling the components in a vertical manner is crucial in order to support modifiability, portability, and reusability. On the other hand components also require some horizontal structuring that is orthogonal to their vertical subdivision.

To achieve these goals, the system is structured into LAYERS so that each layer provides a set of services to the layer above and uses the services of the layer below. Within each LAYER all constituent components work at the same level of abstraction and can interact through connectors. Between two adjacent layers a clearly defined interface is provided. In the pure form of the pattern, layers should not be by-passed: higher-level layers access lower-level layers only through the layer beneath.

Each layer offers a dedicated EXPLICIT INTERFACE [BH03] to the higher-level layers, which remains stable, whereas internal implementation details can change. This way the LAYERS

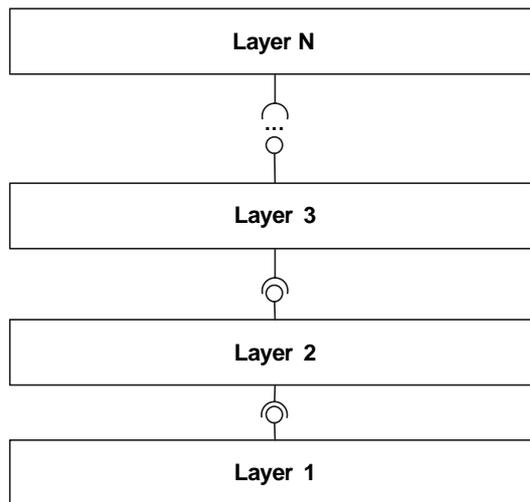


Figure 3: Layers example

pattern allows the work to be sub-divided along clear boundaries, which enables the division of labor. Two adjacent LAYERS can be considered as a CLIENT-SERVER pair, the higher layer being the client and the lower layer being the server. Also, the logic behind layers is especially obvious in the INDIRECTION LAYER where a special layer “hides” the details of a component or subsystem and provides access to its services. An example of LAYERS is shown in Figure 3.

LAYERS is useful for separating higher-level from lower-level responsibilities. On the contrary, the patterns PIPES AND FILTERS and SHARED REPOSITORY place all components at the same level of abstraction. However, both of these patterns may use the LAYERS pattern for structuring the internal architecture of individual architecture elements.

A MICROKERNEL is a layered architecture with three LAYERS: external servers, the microkernel, and internal servers. Similarly the PRESENTATION-ABSTRACTION-CONTROL pattern also enforces LAYERS: a top layer with one agent, several intermediate layers with numerous agents, and one bottom layer which contains the “leaves” agents of the tree-like hierarchy.

A special layered architecture, which allows for implementing many of the following patterns (such as REFLECTION, VIRTUAL MACHINE, and INTERCEPTOR), is INDIRECTION LAYER:

**Pattern: Indirection Layer**

A sub-system should be accessed by one or more components, but direct access to the sub-system is problematic. For instance, the components should not get hard-wired into the sub-system, instead the accessors for the sub-system should be reused. Or the access should be defined in a way that it can be flexibly adapted to changes. The same problem appears at different levels of scale: it can happen between two ordinary components in one environment, components in two different languages, components in two different systems (e.g. if a legacy system is accessed).

An INDIRECTION LAYER is a LAYER between the accessing component and the “instructions” of the sub-system that needs to be accessed. The general term “instructions” can refer to a whole

programming language, or an application programming interface (API) or the public interface(s) of a component or sub-system, or other conventions that accessing components must follow. The INDIRECTION LAYER wraps all accesses to the relevant sub-system and should not be bypassed. The INDIRECTION LAYER can perform additional tasks while deviating invocations to the sub-system, such as converting or tracing the invocations.

The INDIRECTION LAYER can either be integrated to the sub-system (as in “virtual machine” [CBB<sup>+</sup>02]) or be an independent entity (as in ADAPTER or FACADE [GHJV94]) that forwards the invocations to the sub-system. In both cases the accessing components are not aware of this, since the INDIRECTION LAYER aims at exactly that: hiding whatever it is that provides the services.

An example of an INDIRECTION LAYER architecture is shown in Figure 4. It shows a very simple form of INDIRECTION LAYER, consisting of wrappers that are independent of the sub-system providing the services. These wrappers forward the invocations to the hidden components of the subsystem, and perform some actions before and after the invocations. They can also be used to introduce add-on tasks such as logging. More complex INDIRECTION LAYERS, such as VIRTUAL MACHINES or INTERPRETERS, follow the same principle architecture, but perform more complex tasks and hide more components to which they forward requests.

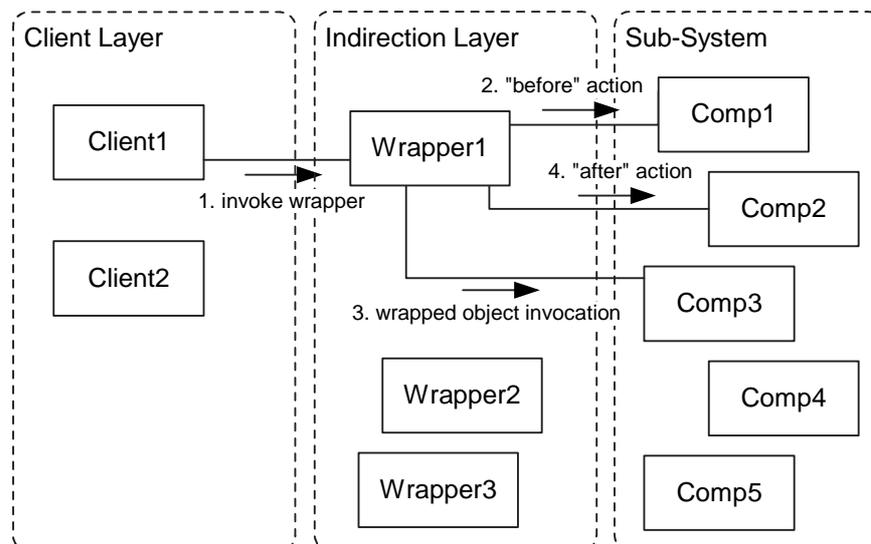


Figure 4: Indirection Layer example: simple wrapper layer

The INDIRECTION LAYER pattern can be thought of as a system of two or three LAYERS: the INDIRECTION LAYER can be considered as Layer N which provides its services to Layer N+1, without disclosing the implementation details of the services to the latter. Sometimes the INDIRECTION LAYER provides all services itself, sometimes the services are provided by yet another Layer N-1 (e.g. a subsystem that is hidden by the INDIRECTION LAYER).

The INDIRECTION LAYER pattern is a foundation for the architectures of INTERPRETER, VIRTUAL MACHINE, and RULE-BASED SYSTEM. These patterns all provide an execution environment for a language defined on some platform. The INTERPRETER, VIRTUAL MACHINE,

or RULE-BASED SYSTEM interpose an INDIRECTION LAYER between the instructions of that language and the instructions of the platform.

The REFLECTION pattern might also be implemented using an INDIRECTION LAYER, so that the latter provides the reflective capabilities of the components defined on top of it. Specifically, the INDIRECTION LAYER can intercept all invocations of the components, and thus can use this information to record the current structure and behavior of the system. This information can then be provided using a reflection API.

## 4 Data Flow View

In the Data Flow View the system is viewed as a number of subsequent transformations upon streams of input data. The concerns addressed by this view are:

- What are the elements that perform the transformations?
- What are the elements that carry the streams of data?
- How are the two aforementioned types of elements connected to each other?
- How are the quality attributes of modifiability, reusability, and integrability supported?

The elements that perform the transformations are components that are independent of one another, and have input and output ports. The elements that carry the streams of data are connectors and similarly have data-in and data-out roles. The relationships between these elements are attachments that connect input ports of components to data-out roles of connectors, and output ports of components to data-in roles of connectors.

### **Pattern: Batch Sequential**

Consider a complex task that can be sub-divided into a number of smaller tasks, which can be defined as a series of independent computations. This should not be realized by one monolithic component because this component would be overly complex, and it would hinder modifiability and reusability.

In a BATCH SEQUENTIAL architecture the whole task is sub-divided into small processing steps, which are realized as separate, independent components. Each step runs to completion and then calls the next sequential step until the whole task is fulfilled. During each step a batch of data is processed and sent as a whole to the next step.

BATCH SEQUENTIAL is a simple sequential data processing architectural pattern. It is useful for simple data flows, but entails severe overhead for starting the batch processes and transmitting data between them. PIPES AND FILTERS is more suitable for stream-oriented data-flow processing, where the filters incrementally transform their input streams into output streams.

### Pattern: Pipes and Filters

Consider as in BATCH SEQUENTIAL the case where a complex task can be sub-divided into a number of smaller tasks, which can be defined as a series of independent computations. Additionally the application processes streams of data, i.e. it transforms input data streams into output data streams. This functionality should not be realized by one monolithic component because this component would be overly complex, and it would hinder modifiability and reusability. Furthermore, different clients require different variations of the computations, for instance, the results should be presented in different ways or different kinds of input data should be provided. To reach this goal, it must be possible to flexibly compose individual sub-tasks according to the client's demands.

In a PIPES AND FILTERS architecture a complex task is divided into several sequential sub-tasks. Each of these sub-tasks is implemented by a separate, independent component, a filter, which handles only this task. Filters have a number of inputs and a number of outputs and they are connected flexibly using pipes but they are never aware of the identity of adjacent filters. Each pipe realizes a stream of data between two components. Each filter consumes and delivers data incrementally, which maximizes the throughput of each individual filter, since filters can potentially work in parallel. Pipes act as data buffers between adjacent filters. The use of PIPES AND FILTERS is advisable when little contextual information needs to be maintained between the filter components and filters retain no state between invocations. PIPES AND FILTERS can be flexibly composed. However, sharing data between these components is expensive or inflexible. There are performance overheads for transferring data in pipes and data transformations, and error handling is rather difficult.

An example of PIPES AND FILTERS is shown in Figure 5. Forks/joins as well as feedback loops are allowed in this pattern, but there is also a variant referred to as a *pipeline*, that forbids both, i.e. has a strict linear topology.

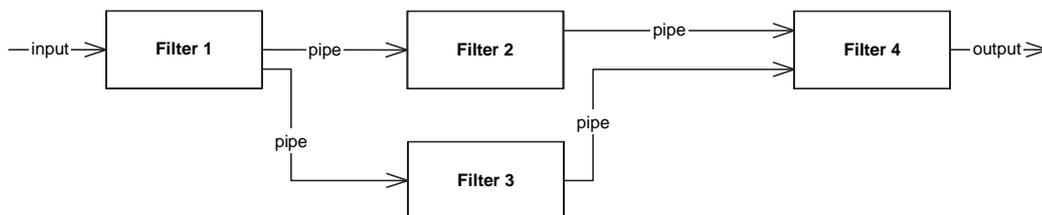


Figure 5: Pipes and filters example

In contrast to BATCH SEQUENTIAL, where there is no explicit abstraction for connectors, the PIPES AND FILTERS pattern considers the pipe connector to be of paramount importance for the transfer of data streams. The key word in PIPES AND FILTERS is flexibility in connecting filters through pipes in order to assemble custom configurations that solve specific problems. Also in PIPES AND FILTERS there is a constant flow of data streams between the filters, while in BATCH SEQUENTIAL, the processing steps are discrete in the sense that each step finishes before the next step may commence.

The pure form of the PIPES AND FILTERS pattern entails that only two adjacent filters can share data through their pipe, but not non-adjacent filters. Therefore pure PIPES AND FILTERS is an alternative to LAYERS and SHARED REPOSITORIES, only if data sharing between non-adjacent processing tasks is not needed. On the other hand, more relaxed forms of the PIPES AND FILTERS pattern can be combined with data-centered architectures like SHARED REPOSITORY, ACTIVE REPOSITORY, or BLACKBOARD to allow for data-sharing between filters. PIPES AND FILTERS can also be used for communication between LAYERS, if data flows through layers are needed.

## 5 Data-centered View

In the Data-centered View the system is viewed as a persistent, shared data store that is accessed and modified by a number of elements. The concerns addressed by this view are:

- How is the shared data store created, accessed, and updated?
- How is data distributed?
- Is the data store passive or active, i.e. does it notify its accessors or are the accessors responsible of finding data of interest to them?
- How does the data store communicate with the elements that access it?
- Do the accessor elements communicate indirectly through the shared data or also directly with each other?
- How are the quality attributes of scalability, modifiability, reusability, and integrability supported?

The data store and the elements that access it are components. The data store is independent of the components, and the components are usually independent of one another. It is possible that there is more than one data store. The elements that transfer data written or read from the data stores are connectors that are attached to the data store(s) and the accessors.

### **Pattern: Shared Repository**

Data needs to be shared between components. In sequential architectures like LAYERS or PIPES AND FILTERS the only way to share data between the components (layers or filters) is to pass the information along with the invocation, which might be inefficient for large data sets. Also it might be inefficient, if the shared information varies from invocation to invocation because the components' interfaces must be prepared to transmit various kinds of data. Finally the long-term persistence of the data requires a centralized data management.

In the SHARED REPOSITORY pattern one component of the system is used as a central data store, accessed by all other independent components. This SHARED REPOSITORY offers suitable means for accessing the data, for instance, a query API or language. The SHARED REPOSITORY must be scalable to meet the clients' requirements, and it must ensure data consistency. It must

handle problems of resource contention, for example by locking accessed data. The SHARED REPOSITORY might also introduce transaction mechanisms.

An example of a SHARED REPOSITORY architecture is shown in Figure 6.

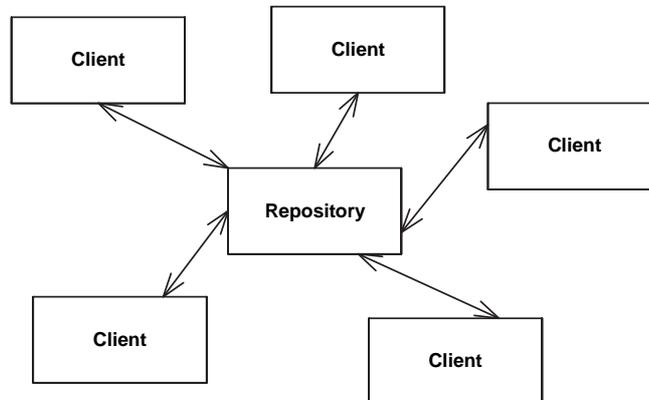


Figure 6: Shared repository example

A SHARED REPOSITORY also might offer additional services, such as security. Some systems offer higher-level access mechanisms, such as query languages or tuple spaces [GCCC85].

A SHARED REPOSITORY offers an alternative to sequential architectures for structuring software components, such as LAYERS and PIPES AND FILTERS, that should be considered, when data sharing or other interaction between non-adjacent components is needed. SHARED REPOSITORIES can be used in a PIPES AND FILTERS architecture to allow for data sharing between filters.

A SHARED REPOSITORY, where all its clients are independent components, can be considered as CLIENT-SERVER, with the data store playing the server part. Similarly it can be considered as a system of two LAYERS where the higher level of clients uses the services of the lower level of the SHARED REPOSITORY.

A variant of the SHARED REPOSITORY pattern is the ACTIVE REPOSITORY pattern<sup>2</sup>:

**Pattern: Active Repository**

A system needs to have a SHARED REPOSITORY, but it should not just be *passively* accessed by accessor components. Clients need to be immediately informed of specific events in the shared repository, such as changes of data or access of data. “Polling” (i.e. querying in frequent intervals) the SHARED REPOSITORY for such events does not work, for instance, because this does not deliver timely information or inflicts overhead on the system performance.

<sup>2</sup>Note that this pattern is introduced as a style named “Blackboard” in [SC97, SG96, BCK98]. We renamed it for this pattern language to ACTIVE REPOSITORY to avoid confusion with the BLACKBOARD pattern from [BMR<sup>+</sup>96] which we discuss below.

An ACTIVE REPOSITORY is a SHARED REPOSITORY that is “active” in the sense that it informs a number of subscribers of specific events that happen in the shared repository. The ACTIVE REPOSITORY maintains a registry of clients and informs them through appropriate notification mechanisms.

The notification mechanism can be realized using ordinary EXPLICIT INVOCATIONS, but in most cases IMPLICIT INVOCATIONS, such as PUBLISH-SUBSCRIBE, are more appropriate.

Another variant of the SHARED REPOSITORY pattern is the BLACKBOARD pattern, which is appropriate when a SHARED REPOSITORY is used in an immature domain in which no deterministic approach to a solution is known or feasible.

**Pattern: Blackboard**

Consider the case where a SHARED REPOSITORY is needed for the shared data of a computation, but no deterministic solution strategies are known. Examples are image recognition or speech recognition applications. However, it should be possible to realize a solution for these types of applications.

In a BLACKBOARD architecture the complex task is divided into smaller sub-tasks for which deterministic solutions are known. The BLACKBOARD is a SHARED REPOSITORY that uses the results of its clients for heuristic computation and step-wise improvement of the solution. Each client can access the BLACKBOARD to see if new inputs are presented for further processing and to deliver results after processing. A control component monitors the blackboard and coordinates the clients according to the state of the blackboard.

An example of a BLACKBOARD architecture is shown in Figure 7. Even though the control component is designed as a separate component, it may as well be part of the clients, the blackboard itself, or a combination of the above.

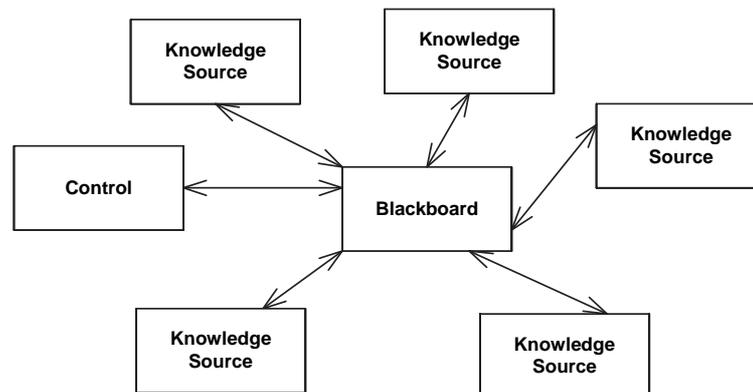


Figure 7: Blackboard example

## 6 Adaptation View

In the Adaptation View the system is viewed as a core part that remains invariable and an adaptable part that either changes over time or in different versions of a system. The concerns addressed by this view are:

- How can a system adapt better to evolution over time or to multiple different versions of a basic architecture?
- What is the system functionality that is more likely to change and what will possibly remain invariable?
- How do the invariable parts communicate with the adaptable parts?
- How are the quality attributes of modifiability, reusability, evolvability, and integrability supported?

The two basic types of elements in this view are the invariable components and the adaptable components (these are often called variation points). These two kinds of components communicate with each other through connectors that have clearly-specified interfaces. Note that some kinds of connectors are adaptable as well, i.e. they are also used as variation points.

Figure 8 summarizes the patterns and their relationships from the Adaptation View and the Language Extension View.

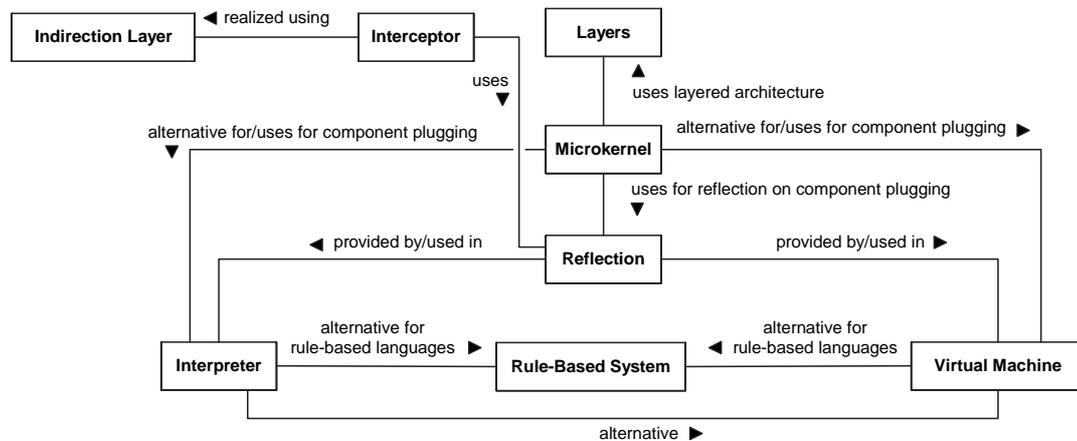


Figure 8: Overview: patterns of the Adaptation and Language Extension View

### Pattern: Microkernel

Consider a system family where different versions of a system need to be supported. In each version, components can be composed in different ways and other details, such as the offered services, public APIs, or user interfaces, might be different. Nonetheless, the system family should be realized using a common architecture to ease software maintenance and foster reuse.

A MICROKERNEL realizes services that all systems, derived from the system family, need and a plug-and-play infrastructure for the system-specific services. Internal servers (not visible to clients) are used to realize version-specific services and they are only accessed through the MICROKERNEL. On the other hand, external servers offer APIs and user interfaces to clients by using the MICROKERNEL. External servers are the only way for clients to access the MICROKERNEL architecture.

The MICROKERNEL pattern promotes flexible architectures which allow systems to adapt successfully to changing system requirements and interfaces. An example of a MICROKERNEL architecture is shown in Figure 9.

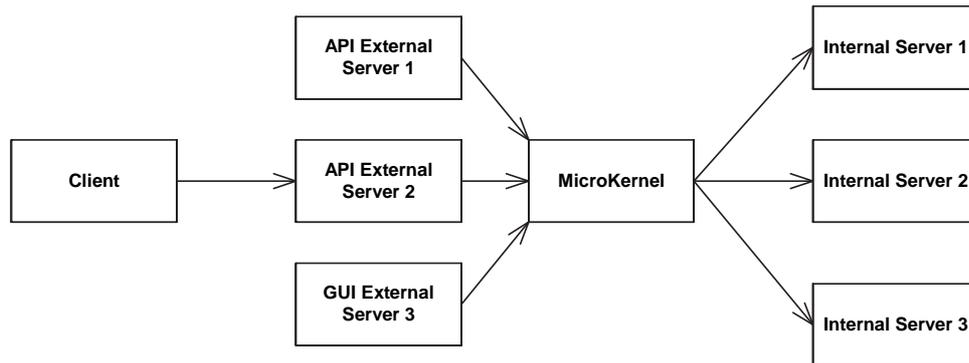


Figure 9: Microkernel example

MICROKERNELS are usually structured in LAYERS: the lowest layer implements an abstraction of the system platform, the next layer implements the services (of the internal servers) used by the MICROKERNEL, the next layer implements the functionality shared by all application versions, and the highest layer glues the external and internal servers together. Apparently, the lowest layer is in fact an INDIRECTION LAYER hiding the low-level system details from the application logic.

In some areas, the patterns INTERPRETER and VIRTUAL MACHINE are alternatives to MICROKERNEL since they all offer a way to integrate or glue components. The MICROKERNEL can integrate version-specific components, while most INTERPRETERS and VIRTUAL MACHINES also offer some way to be extended with components. Thus all three patterns can be used to develop a plug-and-play environment for components. For instance, most scripting languages use their INTERPRETER to offer a gluing framework for components written in the language in which the scripting language itself is implemented (e.g. C, C++, or Java). It is even possible to combine an INTERPRETER or VIRTUAL MACHINE with a MICROKERNEL by implementing the plug-and-play environment of the former as a MICROKERNEL.

A MICROKERNEL introduces an indirection that can be useful in certain CLIENT-SERVER configurations: a client that needs a specific service can request it indirectly through the MICROKERNEL, which establishes the communication to the server that offers this service. In this sense all communication between clients and servers is mediated through the MICROKERNEL, for reasons of e.g. security or modifiability. To develop distributed MICROKERNEL architectures, the

MICROKERNEL can be combined with the BROKER pattern to hide the communication details between clients that request services and servers that implement them.

For all environments that support plug-and-play of components, REFLECTION is useful because it allows to find out which components are currently composed in which way.

**Pattern: Reflection**

Software systems constantly evolve and change over the time, and unanticipated changes are often required. It is hard to automatically cope with changes that are not foreseen.

In a REFLECTION architecture all structural and behavioral aspects of a system are stored into meta-objects and separated from the application logic components. The latter can query the former (that may have changed at any point of time) in order to execute their functionality. Thus REFLECTION allows a system to be defined in a way that allows for coping with unforeseen situations automatically.

The REFLECTION pattern is organized into LAYERS: the meta-level contains the meta-objects which encapsulate the varying structure and behavior; and the base level contains the application logic components that depend on the meta-objects. However this is not the pure LAYERS pattern since not only the base layer uses the services of the meta layer but also the opposite may happen. This is useful for building a reflective system from scratch. The REFLECTION can also be realized with other architectures. For instance, many existing INTERPRETERS and VIRTUAL MACHINES are reflective in the sense that the information in the implementations of the patterns can be used to provide REFLECTION.

An example of a REFLECTION architecture is shown in Figure 10.

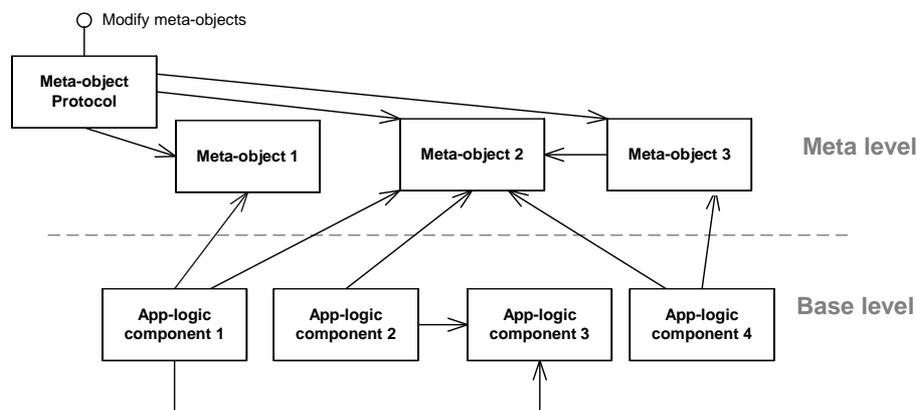


Figure 10: Reflection example

In [Zdu04, Zdu03] more detailed patterns about the realization of REFLECTION in the context of aspect-oriented composition frameworks are presented. The respective variant of REFLECTION is the pattern INTROSPECTION OPTION (see Figure 12).

INDIRECTION LAYER is a more generic pattern than REFLECTION. It can be used to build a REFLECTION-infrastructure, but also other “meta-level” patterns, such as INTERPRETER or

VIRTUAL MACHINE.

In cases where we need an adaptable *framework* to accommodate future services, the INTERCEPTOR pattern is appropriate:

**Pattern: Interceptor**

A framework offers a number of reusable services to the applications that extend it. These services need to be updated in the future as the application domain matures and they should still be offered by the framework, so that the application developers do not need to re-implement them. Furthermore, the framework developer cannot predict all such future services at the point of time where the framework is created, while application developers may not be able to add unanticipated extensions to the framework, in case e.g. that the framework is a black-box.

An INTERCEPTOR is a mechanism for transparently updating the services offered by the framework in response to incoming events. An application can register with the framework any number of INTERCEPTORS that implement new services. The framework facilitates this registration through dispatchers that assign events to INTERCEPTORS. The framework also provides the applications with the means to introspect on the framework's behavior in order to properly handle the events.

The INTERCEPTOR pattern can be realized using an INDIRECTION LAYER or one of its variants, such as INTERPRETER or VIRTUAL MACHINE. The incoming events are therefore re-routed through the INDIRECTION LAYER that consists of several INTERCEPTORS, before they are dispatched to the intended receiver.

INTERCEPTOR can use a REFLECTION mechanism in order to query the framework and retrieve the necessary information to process incoming events.

INTERCEPTOR is defined in special variants: for aspect-oriented composition as MESSAGE INTERCEPTOR in [Zdu04, Zdu03]; for middleware architectures as INVOCATION INTERCEPTOR in [VKZ04]. An example of an INTERCEPTOR architecture is shown in Figure 11. In the figure an invocation is intercepted automatically, and an interceptor manager is invoked, instead of the original target component. Two interceptors are configured for the target component, which are invoked before the invocation. After the interceptors, the target component itself is invoked. When the invocation returns, the interceptors are invoked again, this time in reverse order.

## 7 Language Extension View

In the Language Extension View the system is viewed as a part that is native to the software/hardware environment and another part that is not. The concerns addressed by this view are:

- How can a part of the system that is written in a non-native language be integrated with the software system?
- How can the non-native part be translated into the native environment?

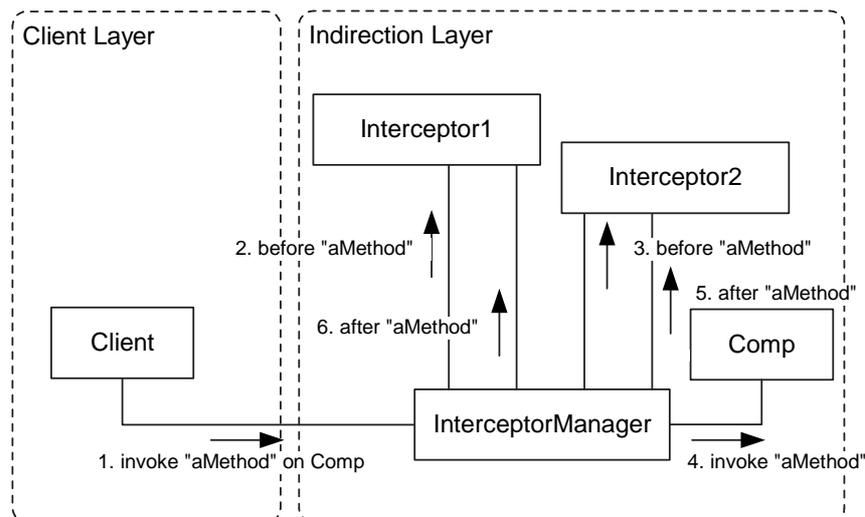


Figure 11: Interceptor example

- How are the quality attribute of portability and modifiability supported?

The native part of the application and the non-native part are components. These communicate indirectly through another type of component, an interpreter component that “translates” the latter into the former. The connectors between these components are data that contain the program instructions in the non-native language, as well as the internal state of the non-native part.

**Pattern: Interpreter**

A language syntax and grammar needs to be parsed and interpreted within an application. The language needs to be interpreted at runtime (i.e. using a compiler is not feasible).

An INTERPRETER for the language is provided, which provides both parsing facilities and an execution environment. The program that needs to be interpreted is provided in form of scripts which are interpreted at runtime. These scripts are portable to each platform realization of the INTERPRETER. For instance, the INTERPRETER can define a class per grammar rule of the language. The parser of the interpreter parses language instructions according to these rules and invokes the interpretation classes. Many more complex INTERPRETER architectures exist.

Some INTERPRETERS use optimizations like on-the-fly byte-code compilers. They thus realize internally elements of a VIRTUAL MACHINE. Note that an INTERPRETER is different to a VIRTUAL MACHINE because it allows for runtime interpretation of scripts, whereas the VIRTUAL MACHINE architecture depends on compilation before runtime:

**Pattern: Virtual Machine**

An efficient execution environment for a programming language is needed. The architecture should facilitate portability, code optimizations, and native machine code generation. Runtime interpretation of the language is not necessarily required.

A VIRTUAL MACHINE defines a simple machine architecture on which not machine code but an intermediate form called the byte-code can be executed. The language is compiled into that byte-code. The VIRTUAL MACHINE can be realized on different platforms, so that the byte-code can be portable between these platforms. The VIRTUAL MACHINE redirects invocations from a byte-code layer into an implementation layer for the commands of the byte-code.

An alternative to INTERPRETERS and VIRTUAL MACHINES, when rule-based or logical languages are needed, is a RULE-BASED SYSTEM:

#### **Pattern: Rule-Based System**

Logical problems are hard to express elegantly in imperative languages that are typically used in INTERPRETERS and VIRTUAL MACHINES. Consider for instance an expert system that provides the knowledge of an expert or a set of constraints. In imperative languages these are expressed by nested if-statements or similar constructs which are rather hard to understand.

A RULE-BASED SYSTEM offers an alternative for expressing such problems in a system. It consists mainly of three things: facts, rules, and an engine that acts on them. Rules represent knowledge in form of a condition and associated actions. Facts represent data. A RULE-BASED SYSTEM applies its rules to the known facts. The actions of a rule might assert new facts, which, in turn, trigger other rules.

As mentioned before INDIRECTION LAYER is the architectural foundation for INTERPRETER, VIRTUAL MACHINE, and RULE-BASED SYSTEM, since either the instructions of the language or the byte-code are re-directed dynamically (at runtime). We do not show component and connector diagrams for these three patterns, since their structure is trivially similar to that of an INDIRECTION LAYER.

In [Zdu04, Zdu03] a number of static alternatives for building language extensions are presented in the context of aspect-oriented composition frameworks. These, for instance, manipulate the parse tree (PARSE TREE INTERPRETER) or byte-code (BYTE-CODE MANIPULATOR) of a language. The relationships to these patterns for aspect-oriented composition are shown in Figure 12.

## **8 User Interaction View**

In the User Interaction View the system is viewed as a part that represents the user interface and a part that contains the application logic, associated with the user interface. The concerns addressed by this view are:

- What is the data and the application logic that is associated to the user interface?
- How is the user interface decoupled from the application logic?
- How are the quality attributes of usability, modifiability, and reusability supported?

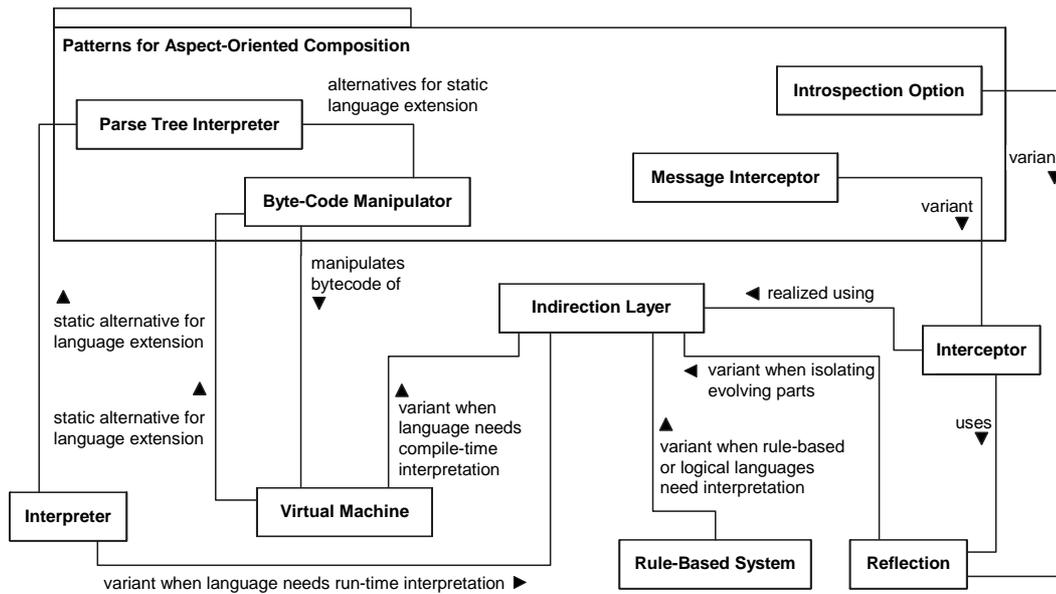


Figure 12: Relationships between patterns in the Language Extension View, Adaptation View and Aspect-oriented Composition View

The elements that present data to the user, accept the user input, and contain the application logic and data are implemented as components. The components interact with each other through connectors that pass data from one to another. This interaction is usually a message-based change notification mechanism.

The relationships between the patterns of this view are shown in Figure 13.

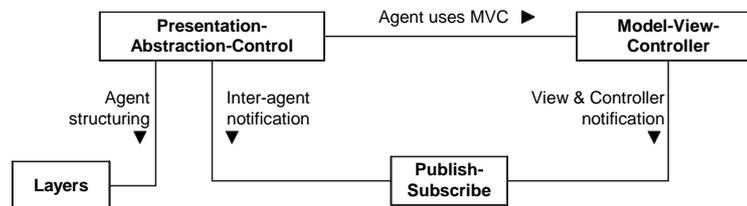


Figure 13: Overview: patterns of the User Interaction View

### Pattern: Model-View-Controller

A system may offer multiple user interfaces. Each user interface depicts all or part of some application data. Changes to the data should be automatically and flexibly reflected to all the different user interfaces. It should be also possible to easily modify any one of the user interfaces, without affecting the application logic associated with the data.

The system is divided into three different parts: a *Model* that encapsulates some application data and the logic that manipulates that data, independently of the user interfaces; one or multiple *Views* that display a specific portion of the data to the user; a *Controller* associated with each *View* that receives user input and translates it into a request to the *Model*. *Views* and *Controllers* constitute the user interface. The users interact strictly through the *Views* and their *Controllers*, independently of the *Model*, which in turn notifies all different user interfaces about updates.

The notification mechanism that updates all *Views* and *Controllers* according to the *Model* can be based on PUBLISH-SUBSCRIBE. All *Controllers* and *Views* subscribe to the *Model*, which in turn publishes the notifications. An example of a MODEL-VIEW-CONTROLLER architecture is shown in Figure 14.

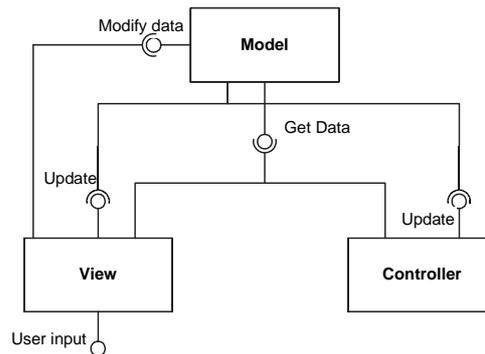


Figure 14: Model-View-Controller example

**Pattern: Presentation-Abstraction-Control**

An interactive system may offer multiple diverse functionalities that need to be presented to the user through a coherent user interface. The various functionalities may require their own custom user interface, and they need to communicate with other functionalities in order to achieve a greater goal. The users need not perceive this diversity but should interact with a simple and consistent interface.

The system is decomposed into a tree-like hierarchy of agents: the leaves of the tree are agents that are responsible for specific functionalities, usually assigned to a specific user interface; at the middle layers there are agents that combine the functionalities of related lower-level agents to offer greater services; at the top of the tree, there is only one agent that orchestrates the middle-layer agents to offer the collective functionality. Each agent is comprised of three parts: a *Presentation* takes care of the user interface; an *Abstraction* maintains application data and the logic that modifies it; a *Control* intermediates between the *Presentation* and the *Abstraction* and handles all communication with the *Controls* of other Agents.

The PRESENTATION-ABSTRACTION-CONTROL pattern is in essence based on MODEL-VIEW-CONTROLLER, in the sense that every agent is designed according to MVC: the *Abstraction* matches the MVC *Model*, while the *presentation* matches the MVC *View* and *Controller*.

On a more macroscopic level, the PRESENTATION-ABSTRACTION-CONTROL patterns is structured according to LAYERS: The top layer contains the chief agent that controls the entire application; the middle layer contains agents with coarse-grained functionality; the lower layer is comprised of fine-grained agents that handle specific services, which users interact with. An example of a PRESENTATION-ABSTRACTION-CONTROL architecture is shown in Figure 15.

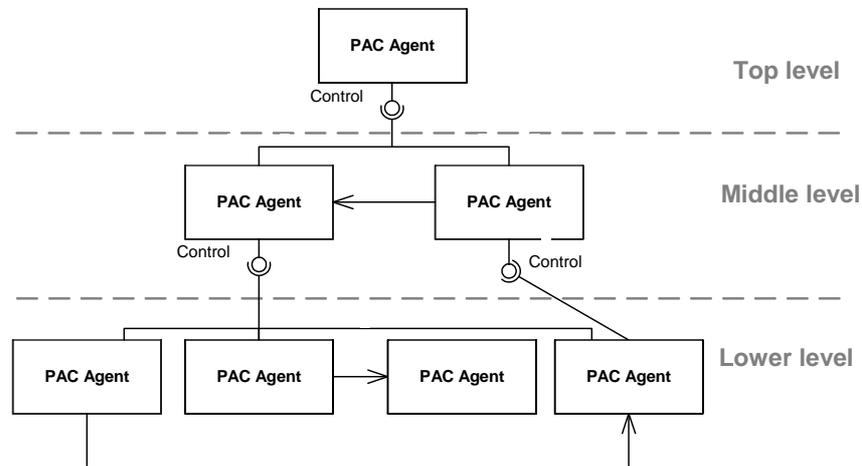


Figure 15: Presentation-Abstraction-Control example

The various agents usually need to propagate changes to the rest of the agent hierarchy, and this can be achieved again through the PUBLISH-SUBSCRIBE pattern. Usually higher-level agents subscribe to the notifications of lower-level agents.

An alternative to MVC and PAC for applications with extensive user interface requirements and other particular requirements is the C2 architectural pattern.

**Pattern: C2**

An interactive system is comprised of multiple components such as GUI widgets, conceptual models of those widgets at various levels, data structures, renderers, and of course application logic. The system may need to support several requirements such as: different implementation language of components, different GUI frameworks reused, distribution in a heterogeneous network, concurrent interaction of components without shared address spaces, run-time reconfiguration, multi-user interaction. Yet the system needs to be designed to achieve separation of concerns and satisfy its performance constraints.

The system is decomposed into a top-to-bottom hierarchy of concurrent components that interact asynchronously by sending messages through explicit connectors. Components submit request messages upwards in the hierarchy, knowing the components above, but they send notification messages downwards in the hierarchy, without knowing the components lying beneath. Components are only connected with connectors, but connectors may be connected to both components and other connectors. The purposes of connectors is to broadcast, route, and filter messages.

An example of a C2 architecture with four components in three layers, and two connectors that delimit the layers, is shown in Figure 16.

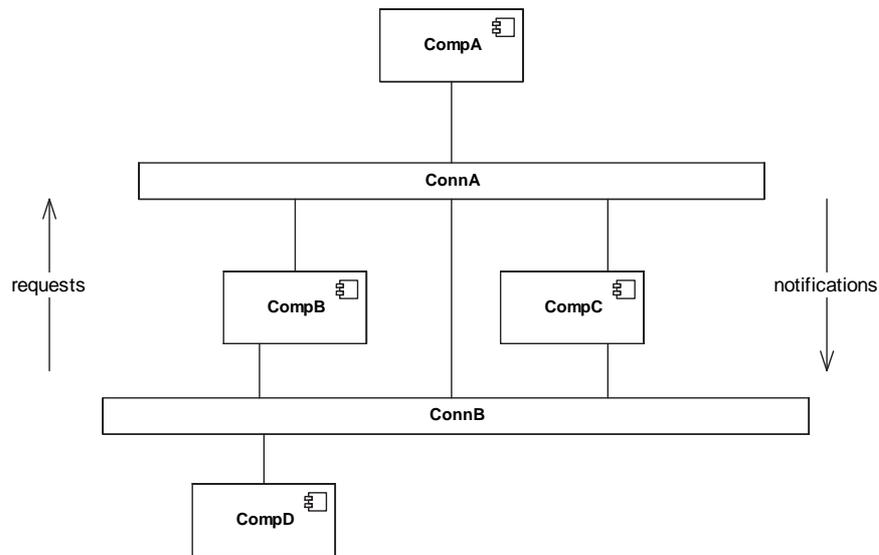


Figure 16: C2 example

The C2 top-to-bottom hierarchy resembles the form of a LAYERS architecture in an upside-down order. A C2 component that belongs to a given layer, uses the services of the layers above it by invoking services on them and provides services to the layers below it by sending notifications to them.

Since the C2 pattern provides *substrate independence* [TMA<sup>+</sup>96], isolating a component from the components underneath it, the layer where a component is placed is in essence an INDIRECTION LAYER.

The interaction between the C2 components takes place through asynchronous message exchange, thus utilizing an IMPLICIT INVOCATION mechanism, and specifically callbacks, e.g. PUBLISH-SUBSCRIBE.

## 9 Component Interaction View

In the Component Interaction View the system is viewed as a number of independent components that interact with each other in the context of a system. The concerns addressed by this view are:

- How do the independent components interact with each other?
- How are the individual components decoupled from each other?
- How are the quality attributes of modifiability and integrability supported?

The components retain their independence, since they merely exchange data but do not directly control each other. The components interact with each other through connectors that pass data from one to another. This interaction can be performed synchronously or asynchronously and can be message-based or through direct calls. This view is closely connected to the distributed view, since the independent components might be distributed in various network nodes or processes.

The two major patterns in this view differentiate whether the components interact through explicit or implicit invocations<sup>3</sup>:

**Pattern: Explicit Invocation**

Consider a component, the client, which needs to invoke a service defined in another component, the supplier. Coupling the client with the supplier in various ways is not only harmless but often desirable. For example the client must know the exact network location of the component which offers the service in order to improve performance; or the client must always initiate the invocation itself; or the client must block, waiting for the result of the invocation, before proceeding with its business; or the topology of the interacting clients and suppliers is known beforehand and must remain fixed. How can these two components interact?

An EXPLICIT INVOCATION allows a client to invoke services on a supplier, by coupling them in various respects. The decisions that concern the coupling (e.g. network location of the supplier) are known at design-time. The client provides these design decisions together with the service name and parameters to the EXPLICIT INVOCATION mechanism, when initiating the invocation. The EXPLICIT INVOCATION mechanism performs the invocations and delivers the result to the client as soon as it is computed. The EXPLICIT INVOCATION mechanism may be part of the client and the server or may exist as an independent component.

An example of an EXPLICIT INVOCATION architecture is shown in Figure 17, where the EXPLICIT INVOCATION mechanism is implemented with the help of a BROKER and a PROXY, as part of both the client and the supplier.

During the EXPLICIT INVOCATION the identification of the service supplier, can be realized, for instance by using the pattern OBJECT ID [VKZ04]. The client also knows the location of the service supplier, and furthermore, in some systems, the service supplier needs to know about the location of the client, so that the result can be sent back. This can be achieved by OBJECT IDS enriched with location information, as mandated by the pattern ABSOLUTE OBJECT REFERENCE [VKZ04].

There are two main variants of EXPLICIT INVOCATIONS: synchronous, explicit invocations and asynchronous, explicit invocations. In a synchronous invocation, the client blocks until the result is available. In an asynchronous invocation, the client continues with its work immediately, and the result is delivered at a later point, after it is computed. There are four patterns that describe different variants of asynchronous invocations for distributed systems [VKZ04]:

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<sup>3</sup>Note that in [BCK98] an EXPLICIT INVOCATION is given the opposite meaning as it is considered a sub-pattern of the EVENT SYSTEMS pattern. However we have chosen this name to show the contrast between components explicitly (e.g. direct method call) and implicitly (e.g. events) invoking each other.



should suspend the invocation until the supplier is up again; or the client and the supplier are part of dissimilar systems and thus the invocation must be transformed, queued, or otherwise manipulated during delivery. How can such additional requirements during delivery be met?

In the IMPLICIT INVOCATION pattern the invocation is not performed explicitly from client to supplier, but indirectly and rather randomly through a special mechanism such as PUBLISH-SUBSCRIBE, MESSAGE QUEUING, or broadcast, that decouples clients from suppliers. All additional requirements for invocation delivery are handled by the IMPLICIT INVOCATION mechanism during the delivery of the invocation.

An example of implicit invocation is the synchronization between Model, View, and Controller in the MODEL-VIEW-CONTROLLER pattern, as depicted in Figure 15. The Model notifies its Views and Controllers whenever its data have been changed, and so Views and Controllers implicitly invoke the Model to get the updated data. The Views and Controllers are decoupled from the Model, since they do not initiate the invocation, but the Model does it when it accepts an certain event. Models and Controllers may also be added and removed dynamically.

IMPLICIT INVOCATION can be both synchronous and asynchronous as can EXPLICIT INVOCATION, meaning that the client can either block or not, waiting for the invocation result. However IMPLICIT INVOCATIONS are most often asynchronous, in contrast to EXPLICIT INVOCATIONS, which are usually synchronous. Thus the aforementioned patterns for asynchronous result handling should be used for IMPLICIT INVOCATION as well. An even more prominent contrast between them is that in EXPLICIT INVOCATION the invocation is always deterministic from client to supplier, while in IMPLICIT INVOCATION the trigger happens randomly (e.g. through an event) and not necessarily initiated by a client (e.g. by the producer in PUBLISH-SUBSCRIBE).

Same as in EXPLICIT INVOCATION, distributed IMPLICIT INVOCATION usually uses a BROKER to hide the details of network communication and allow the components to contain only their application logic. Note that an IMPLICIT INVOCATION mechanism decouples clients from suppliers, while the BROKER pattern decouples both from the communication infrastructure.

There are different IMPLICIT INVOCATION variants, with respect to the tasks performed during the delivery of the invocation. For instance, in a broadcast mechanism the location of the invocation receiver is unknown to the client, since the invocation is broadcast through the network. This variant is used, e.g. for looking up the initial reference in a PEER-TO-PEER system. An event system realizes PUBLISH-SUBSCRIBE, in order to decouple producers and consumers of data. The MESSAGE QUEUING pattern queues invocations and results to increase delivery reliability, handle temporal outages of the supplier, and perform other tasks.

Among the implicit and explicit invocation patterns and their variants, only the synchronous variant of EXPLICIT INVOCATION can align a result unambiguously to an invocation, because the client blocks on the result. For all other cases – when invocations are performed asynchronously, it is possible that one client sends multiple invocations after another, and results for these invocations arrive in a different order than the invocations. Because the same client performs the invocations, the OBJECT ID of the client cannot be used for aligning a result to an invocation. An ASYNCHRONOUS COMPLETION TOKEN [SSRB00] contains information that identifies the individual invocation and perhaps also other information such as a behavior to be executed in the client when the result is processed. The ASYNCHRONOUS COMPLETION TO-

KEN is sent along with each asynchronous invocation of a client, and the service supplier sends it back with the result. Thus the client can use this information to align the result to the invocation. The ASYNCHRONOUS COMPLETION TOKEN is used in IMPLICIT INVOCATIONS and asynchronous EXPLICIT INVOCATION to align invocations to incoming results.

Figure 18 shows the relations of IMPLICIT INVOCATION and EXPLICIT INVOCATION, while Figure 19 gives an overview of all the patterns for component interaction as well as distribution.

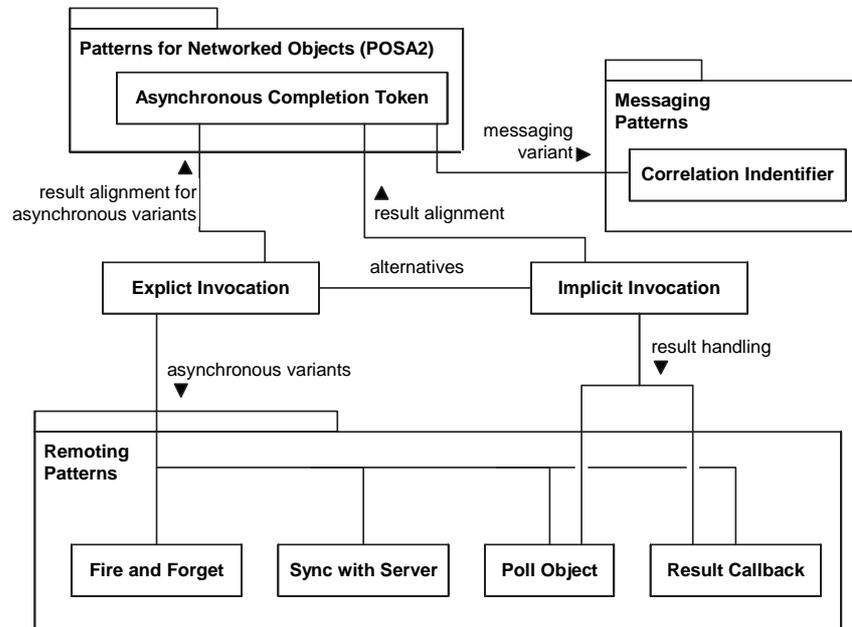


Figure 18: Overview: patterns for basic component interaction

There are two variants of the EXPLICIT INVOCATION pattern: CLIENT-SERVER and PEER-TO-PEER.

**Pattern: Client-Server**

Two components need to communicate, and they are independent of each other, even running in different processes or being distributed in different machines. The two components are not equal peers communicating with each other, but one of them is initiating the communication, asking for a service that the other provides. Furthermore, multiple components might request the same service provided by a single component. Thus, the component providing a service must be able to cope with numerous requests at any time, i.e. the component must scale well). On the other hand, the requesting components using one and the same service might deal differently with the results. This asymmetry between the components should be reflected in the architecture for the optimization of quality attributes such as performance, shared use of resources, and memory consumption.

The CLIENT-SERVER pattern distinguishes two kinds of components: clients and servers. The client requests information or services from a server. To do so it needs to know how to access the

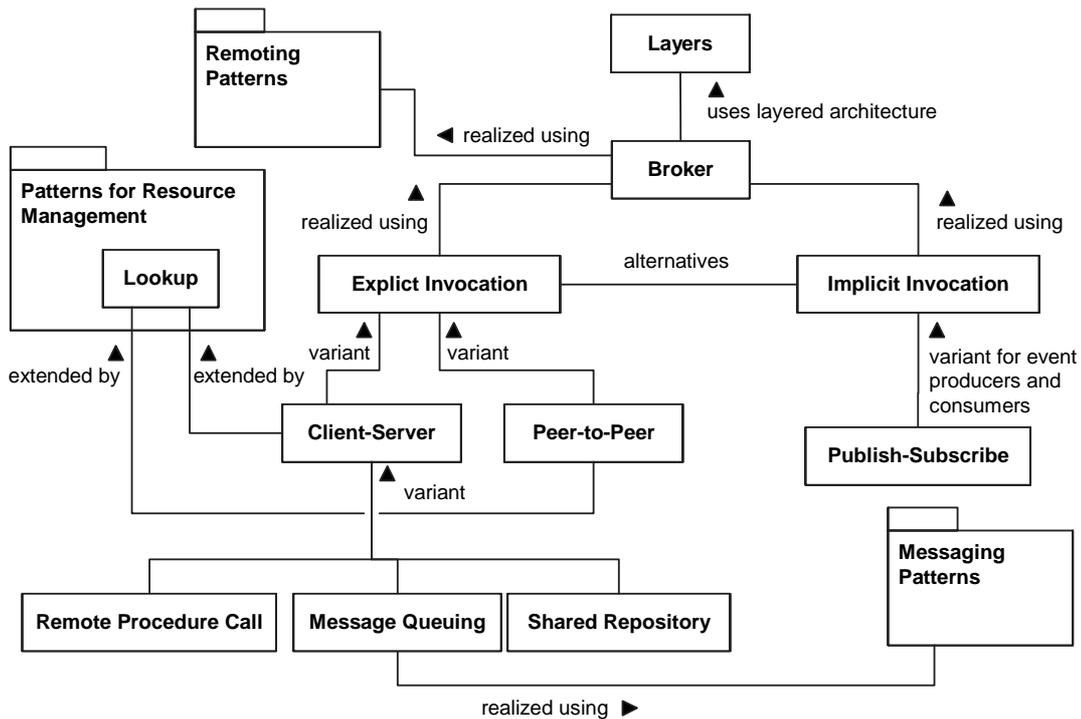


Figure 19: Overview: patterns for component interaction and distribution

server, that is, it requires an ID or an address of the server and of course the server's interface. The server responds to the requests of the client, and processes each client request on its own. It does not know about the ID or address of the client before the interaction takes place. Clients are optimized for their application task, whereas servers are optimized for serving multiple clients.

Both client and server must implement collective tasks, such as security, transaction, and systems management – something that is more complex in a CLIENT-SERVER architecture than in simple EXPLICIT INVOCATIONS.

Sophisticated, distributed CLIENT-SERVER architectures usually rely on the BROKER pattern to make the complexity of the distributed communication manageable. The same is true for the PEER-TO-PEER pattern.

Using the CLIENT-SERVER pattern we can build arbitrarily complex architectures by introducing multiple client-server relationships: a server can act itself as a client to other servers. The result is a so-called *n-Tier-architecture*. A prominent example of such architectures is the 3-tier-architecture (see Figure 20), which consists of:

- a client tier, responsible for the presentation of data, receiving user events, and controlling the user interface
- an application logic tier, responsible for implementing the application logic (also known as business logic)

- a backend tier, responsible for providing backend services, such as data storage in a data base or access to a legacy system

It is also possible to combine the CLIENT-SERVER pattern with LAYERS in order to design a system where the client and the server components individually are layered. For example the ISO/OSI standard defines such an architecture, where there are seven layers in both the client and the server side, and each client communicates with the server at the same layer addressing a certain scope and responsibilities.

Furthermore SHARED REPOSITORIES or BLACKBOARDS can be perceived as CLIENT-SERVER, where the data store is the server and the data accessors are the clients.

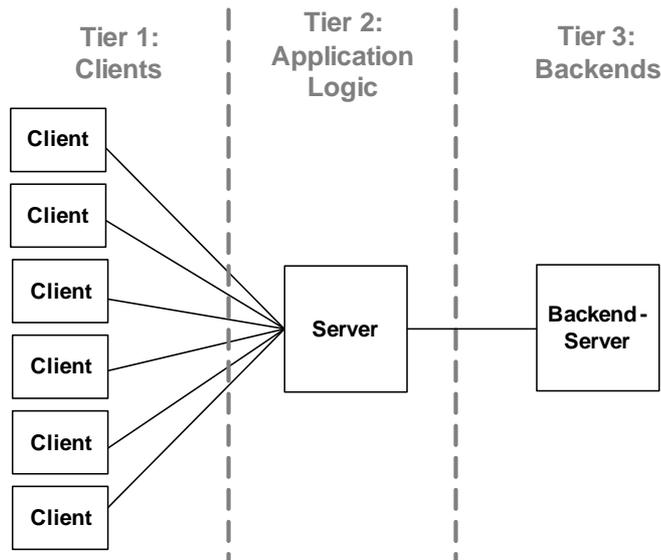


Figure 20: 3-tier client-server architecture: example

When CLIENT-SERVER is used in a distributed fashion, it can be extended so that the location of a remote component does not need to be hard-wired into the system. The pattern LOOKUP [KJ04, VKZ04] allows servers to register their remote components at a central service e.g. by name or property. Using the pattern LOOKUP, the client must only know the location of the lookup service instead of the potentially huge number of locations of the remote components it wants to communicate with. The LOOKUP pattern is thus an alternative to broadcast messages for getting initial references (e.g. using IMPLICIT INVOCATIONS). The problem that remains however is how to get the initial reference to the lookup component.

A general alternative to CLIENT-SERVER is PEER-TO-PEER:

**Pattern: Peer-to-Peer**

Consider a situation similar to that of a CLIENT-SERVER, but in contrast to CLIENT-SERVER, there is no distinction between the components: each component might both provide services and consume services. When a component provides a service it must perform well according

to the demands of the requesting components. Each component must know how to access other components.

In the PEER-TO-PEER pattern each component has equal responsibilities, in particular it may act both as a client and as a server. Each component offers its own services (or data) and is able to access the services in other components. The PEER-TO-PEER network consists of a dynamic number of components. A PEER-TO-PEER component knows how to access the network. Before a component can join a network, it must get an initial reference to this network. This is solved by a bootstrapping mechanism, such as providing public lists of dedicated peers or broadcast messages (using IMPLICIT INVOCATION) in the network announcing peers.

Once an initial reference of the PEER-TO-PEER network is found, we need to find other peers in the network. For this purpose, each peer (or each dedicated peer) realizes the LOOKUP pattern [KJ04, VKZ04]. Using LOOKUP peers can be found based on their names or their properties. PEER-TO-PEER can be realized internally using CLIENT-SERVER, or other patterns. It usually also uses a BROKER architecture.

Whereas CLIENT-SERVER and PEER-TO-PEER concentrate on EXPLICIT INVOCATIONS, PUBLISH-SUBSCRIBE<sup>4</sup> is an interaction pattern that is heavily based on IMPLICIT INVOCATIONS:

**Pattern: Publish-Subscribe**

A component should be accessed or informed of a specific runtime event. Events are of different nature than direct interactions as in CLIENT-SERVER or PEER-TO-PEER. Sometimes a number of components should be actively informed (an announcement or broadcast), in other cases only one specific component is interested in the event. In contrast to EXPLICIT INVOCATIONS, event producers and consumers need to be decoupled for a number of reasons: to support locality of changes; to locate them in different processes or machines; to allow for an arbitrary time period between the announcement of interest in an event, and the actual triggering of the event. Still, there must be a way to inform the interested components?

PUBLISH-SUBSCRIBE allows event consumers (subscribers) to register for specific events, and event producers to publish (raise) specific events that reach a specified number of consumers. The PUBLISH-SUBSCRIBE mechanism is triggered by the event producers and automatically executes a callback-operation to the event consumers. The mechanism thus takes care of decoupling producers and consumers by transmitting events between them.

An example of a PUBLISH-SUBSCRIBE architecture, where the PUBLISH-SUBSCRIBE mechanism is implemented as an independent subscription manager, is shown in Figure 21.

In the local context the PUBLISH-SUBSCRIBE can be based on the OBSERVER pattern [GHJV94], where the PUBLISH-SUBSCRIBE mechanism is implemented as part of the ‘subject’ (i.e. the event producer). For instance, most GUI frameworks are based on a PUBLISH-SUBSCRIBE model.

In the remote context PUBLISH-SUBSCRIBE is used in MESSAGE QUEUING implementations or as a remoting pattern of its own. PUBLISH-SUBSCRIBE makes no assumption about the order of

<sup>4</sup>In [BMR<sup>+</sup>96] this pattern is called PUBLISHER-SUBSCRIBER and in [BCK98] it is called EVENT SYSTEM.

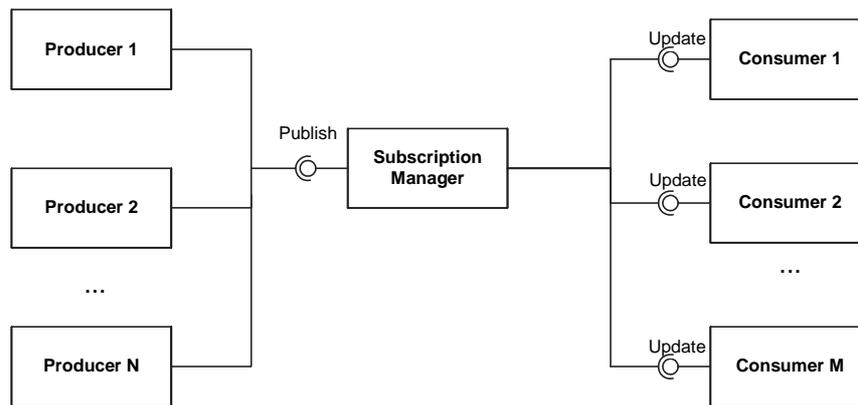


Figure 21: Publish-Subscribe example

processing events. This and other issues are solved by some messaging patterns [HW03] (see MESSAGE QUEUING explained below).

The PUBLISH-SUBSCRIBE pattern can be used in the context of the ACTIVE REPOSITORY pattern, so that accessors of data subscribe to the repository, which in turn notifies them when the data is updated.

PUBLISH-SUBSCRIBE is sometimes used to realize CLIENT-SERVER and PEER-TO-PEER: for instance, in distributed implementations of CLIENT-SERVER and PEER-TO-PEER it is necessary to bridge between the asynchronous network events and the synchronous processing model of the server. This can be done using a local PUBLISH-SUBSCRIBE model, where event handlers subscribe for the network events. This is described in detail in the REACTOR pattern [SSRB00].

In a similar way, PUBLISH-SUBSCRIBE models are used for integration tools, to ensure database consistency, and in user interfaces. For example, in the MODEL-VIEW-CONTROLLER pattern, the view and controllers subscribe to the model which then publishes any updates to them. Similarly in the PRESENTATION-ABSTRACTION-CONTROL pattern, higher-level agents subscribe to lower-level agents that handle the user interface. In general, PUBLISH-SUBSCRIBE offers a high potential for reuse and evolution, because it offers a strong decoupling between event producers and event consumers. There are also some potential problems, however: event consumers have to register for events, which is in general more complex than for instance a CLIENT-SERVER interaction. There is no guarantee that an event will be processed. Exchange of data is not as simple as parameter passing. Often a SHARED REPOSITORY must be used, which is slower than parameter passing.

## 10 Distribution view

In the Distribution View the system is viewed as a number of components that are distributed among network nodes (or different processes). The concerns addressed by this view are:

- How do the distributed components interact with each other?

- How are the distributed components decoupled from each other?
- How are the quality attributes of interoperability, location-transparency, performance, and modifiability supported?

The components are physically located in different network nodes or processes. They interact with each other through connectors that pass invocations or data from one to another. The details of these interactions can be better explained in the Component Interaction View.

**Pattern: Broker**

Distributed software system developers face many challenges that do not arise in single-process software. One is the communication across unreliable networks. Others are the integration of heterogeneous components into coherent applications, as well as the efficient use of networking resources. If developers of distributed systems must overcome all these challenges within their application code, they may lose their primary focus: to develop applications that efficiently tackle their domain-specific problems.

A BROKER separates the communication functionality of a distributed system from its application functionality. The BROKER hides and mediates all communication between the objects or components of a system. A BROKER consists of a client-side REQUESTOR [VKZ04] to construct and forward invocations, as well as a server-side INVOKER [VKZ04] that is responsible for invoking the operations of the target remote object. A MARSHALLER [VKZ04] on each side of the communication path handles the transformation of requests and replies from programming-language native data types into a byte array that can be sent over the transmission medium.

The BROKER is a compound pattern that is realized using a number of remoting patterns [VKZ04]. The most foundational remoting patterns in a BROKER architecture are mentioned above: REQUESTOR, INVOKER, and MARSHALLER. There are many others. Some important examples are: a CLIENT PROXY [VKZ04] represents the remote object in the client process. This proxy has the same interface as the remote object it represents. An INTERFACE DESCRIPTION [VKZ04] is used to make the remote object's interface known to the clients. LOOKUP [KJ04, VKZ04] allows clients to discover remote objects.

The BROKER uses a LAYERS architecture. The layers of BROKER are also described in [VKZ04].

Many well-known BROKER realizations are based on the CLIENT SERVER pattern. However, the other patterns for component interactions, such as EXPLICIT INVOCATION, PEER-TO-PEER, MESSAGE QUEUING, and PUBLISH-SUBSCRIBE, can also use a BROKER to isolate communication-related concerns, when used in a distributed setting. Only in very simple distributed systems or in distributed systems with severe constraints (e.g. regarding performance and memory consumption), it might be advisable not to use a BROKER.

An example for a BROKER realization is depicted in Figure 22.

The BROKER can be seen as the general structure that utilizes the patterns from the Component Interaction View in a distributed setting. The following remoting patterns [VKZ04] are a number of variants of CLIENT-SERVER that usually operate in a distributed setting and are mutual alternatives. They usually employ a BROKER architecture internally.

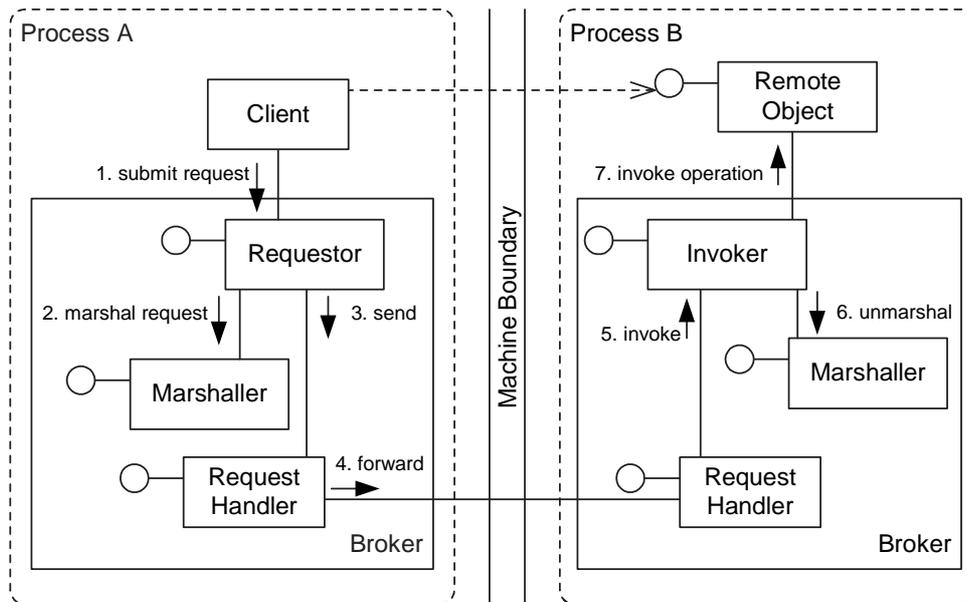


Figure 22: Example of a Broker architecture

#### Pattern: Remote Procedure Calls

Consider the case where you want to realize an EXPLICIT INVOCATION in a distributed setting. The use of low-level network protocols requires developers to invoke the send and receive operations of the respective network protocol implementations. This is undesirable because the network access code cannot be reused, low-level details are not hidden, and thus solutions are hard to maintain and understand.

REMOTE PROCEDURE CALLS extend the well-known procedure call abstraction to distributed systems. They aim at letting a remote procedure invocation behave as if it were a local invocation. Programs are allowed to invoke procedures (or operations) in a different process and/or on a remote machine.

REMOTE PROCEDURE CALLS leverage the CLIENT-SERVER pattern of interaction: a client invokes operations, and a server provides a well-defined set of operations that the client can invoke. To the client developer, these operations look almost like local operations. A major difference is that additional errors might occur during a remote invocation, for example because the network fails or the requested operation is not implemented by the server. These errors must be signaled to the client (see the pattern REMOTING ERROR [VKZ04]).

#### Pattern: Message Queuing

Consider a situation similar to that of REMOTE PROCEDURE CALLS, but it is necessary to decouple the sender from the receiver to realize queuing of invocations. Queuing is necessary, for instance, when temporal outages of the receiver should be tolerated or when heterogeneous systems should be integrated. For instance, when a legacy system using BATCH SEQUENTIAL

should be integrated into a distributed system, only one invocation can be handled at a time by that system. Somewhere additional messages must be queued until the system is ready to process the next message.

Messages are not passed from client to server application directly, but through intermediate message queues that store and forward the messages. This has a number of consequences: senders and receivers of messages are decoupled, so they do not need to know each other's location (perhaps not even the identity). A sender just puts messages into a particular queue and does not necessarily know who consumes the messages. For example, a message might be consumed by more than one receiver. Receivers consume messages by monitoring queues.

MESSAGE QUEUING realizes CLIENT-SERVER interactions and implements IMPLICIT INVOCATION as the primary invocation pattern.

MESSAGE QUEUING architectures are explained in detail in [HW03] using a pattern language for messaging systems. An example of a MESSAGE QUEUING architecture is shown in Figure 23.

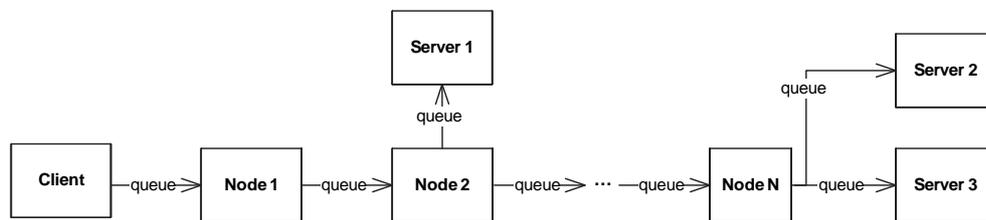


Figure 23: Message Queuing example

## 11 Epilogue

We have proposed to unite existing approaches of architectural patterns into a pattern language so that practitioners can benefit from a single comprehensive source of patterns. We emphasized on outlining the relations between the architectural patterns in order for them to acquire added value as a language, rather than a set of individual patterns. We also referenced the original sources for the patterns, so that interested parties can explore the rich details of each pattern.

Organizing the entire set of architectural patterns into a coherent pattern language is an immense amount of work; therefore we limited this paper to including patterns from the initial and fundamental catalogues and categorizations that deal with components and connectors. We hope that other members of the community can give us feedback so that we can at least reach a consensus on these fundamental patterns.

## Acknowledgments

We like to thank our EuroPLoP 2005 shepherd Lars Grunske.

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